

## ON THE BOUNDED QUASI-DEGREES OF C.E. SETS

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Tennenbaum (see [1,p.159]) defined the notion of Q-reducibility on sets of natural numbers as follows: a set A is Q-reducible to a set B (in symbols:  $A \leq_Q B$ ) if there exists a computable function f such that for every  $x \in \omega$  (where  $\omega$  denotes the set of natural numbers),

$$x \in A \Leftrightarrow W_{f(x)} \subseteq B.$$

In this case we say that  $A \leq_Q B$  via f.

If  $A \leq_Q B$  via f and there exists a fixed number  $n \in \omega$  such that for all x,  $|W_{f(x)}| \leq n$ , then we say that the set A is bounded Q-reducible to the set B, denoted by  $A \leq_{bQ} B$ . A computably enumerable (c.e.) set A is bQ-complete if every c.e. set is bQ-reducible to A. The relation of bQ-reducibility is reflexive and transitive so that it generates a degree structure on the subsets of  $\omega$ .

Our notation and terminology are standard and can be found in [1].

**Theorem 1.** For every noncomputable c.e. incomplete bQ-degree, there exists a nonspeedable bQ-degree incomparable with it.

**Corollary .** Among all nonspeedable bQ-degrees contained in a c.e. complete T-degree there is no maximal bQ-degree.

A noncomputable set X is called computable separable (shortly: r-separable), if for every c.e. set Y

$$Y \cap X = \emptyset \Rightarrow (\exists R \text{ computable}) (X \subseteq R \& Y \cap R = \emptyset).$$

**Theorem 2.** Let A and B be c.e. sets and suppose that A is an r-separable set. Then if  $A \leq_{bQ} B$ , then there exists a noncomputable c.e. C such that  $C \leq_m A$  &  $C \leq_m B$ .

**Corollary.** Let M be a maximal set and A a c.e. set. Then  $M \leq_{bQ} A \Rightarrow M \leq_m A$ .

**Theorem 3.** The structure  $\mathcal{D}_{bs}$  of the bs-degrees is not elementary equivalent neither to the structure  $\mathcal{D}_{be}$  of the be-degrees nor to the structure  $\mathcal{D}_e$  of the e-degrees.

**Corollary.** The structure  $\mathcal{D}_{bQ}$  of the bQ-degrees is not elementary equivalent neither to the structure  $\mathcal{D}_{be}$  of the be-degrees nor to the structure  $\mathcal{D}_e$  of the e-degrees.

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[1] H.Rogers, Theory of recursive functions and effective computability. McGraw-Hill Book Co., New York, 1967.